

REPEATED BURST ERROR CORRECTING LINEAR CODES

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Many kinds of errors in coding theory have been dealt with for which codes have been constructed to combat such errors. Though there is a long history concerning the growth of the subject and many of the codes developed have found applications in numerous areas of practical interest, one of the areas of practical importance in which a parallel growth of the subject took place is that of burst error detecting and correcting codes. The nature of burst errors differ from channel to channel depending upon the behaviour of channels or the kind of errors which occur during the process of data transmission. In very busy communication channels, errors repeat themselves more frequently. In view of this, it is desirable to consider repeated burst errors. The paper presents lower and upper bounds on the number of parity-check digits required for a linear code correcting errors in the form of repeated bursts. An upper bound for a code that detects m -repeated bursts has also been derived. Illustrations of several codes that correct 2-repeated bursts of different lengths have also been given.

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1. Introduction

In the theory of error correcting codes, there are many kinds of errors for which codes have been developed to detect and correct such errors. Codes developed at the early stages were meant mainly to detect and correct random errors. However, it was noticed later that in many channels, the likelihood of the occurrence of errors was more in adjacent positions rather than their occurrence in a random manner. It was in this spirit that codes correcting single errors and double adjacent errors were

developed by Abramson [1]. Upon generalization of this idea, such errors were put in the category of errors called burst errors. It has been noticed that in many channels, errors occur predominantly in the form of bursts. Since the development of various burst error correcting and detecting codes, several variants and modifications of the definition of a burst error came up depending upon the various kinds of channels which were in use. A burst of length b may be defined as follows:

Definition 1.1. A burst of length b is a vector whose only non-zero components are among some b consecutive components, the first and the last component of which is non-zero.

This definition was given by P. Fire [6] in a research report wherein he called such errors as ‘open-loop burst errors’ and in the same report he considered another kind of burst errors as well which were called ‘closed-loop burst errors’. A closed-loop burst of length b is defined as follows:

Definition 1.2. Let b be an integer and $x = (\xi_1, \xi_2, \dots, \xi_n)$ be a vector in $V^n(q)$, a vector space of n -tuples over $\text{GF}(q)$. If $2 \leq b \leq \frac{(n+1)}{2}$, then x is called a ‘closed-loop burst vector of length b ’ whenever there is an i such that $1 \leq i \leq b-1, \xi_i \cdot \xi_{n-b+i+1} \neq 0, \xi_{i+1} = \xi_{i+2} = \dots = \xi_{n-b+i} = 0$.

There is yet another kind of burst error due to Chien and Tang [4] which has drawn attention of many researchers as follows:

Definition 1.3. A burst of length b is a vector whose only non-zero components are confined to some b consecutive components, the first of which is non-zero.

Still amongst the various generalizations of bursts, there are codes which have been developed for the correction of multiple bursts, bursts of bursts, random multiple bursts, low-density bursts etc. A good deal of research work has been devoted for the development of burst/multiple burst correcting codes (see [15], [14], [8], [12], [3], [16] and others). Amongst the category of multiple bursts, there is no uniform terminology. Looking at the fact that there are many channels which require development of codes dealing with multiple bursts (the channels studied in [10] fall in this category), we in this paper propose a new kind of burst error which will be called repeated burst error. Such an error may be defined as follows:

Definition 1.4. An m -repeated burst of length b is a vector of length n whose only non-zero components are confined to m distinct sets of b consecutive components, the first and the last component of each set being non-zero.

The development of such codes detecting and correcting such errors may prove to be useful for channels already dealing with multiple burst errors thereby improving upon their efficiency. Such codes emerge as a natural generalization of single burst error correcting and detecting codes.

In particular, a 2-repeated burst may be defined as follows:

Definition 1.5. A 2-repeated burst of length b is a vector of length n whose only non-zero components are confined to two distinct sets of b consecutive components, the first and the last components of each set being non-zero.

As an illustration (0012040003012000) is a 2-repeated burst of length 4 over $\text{GF}(5)$.

Consider a 2-repeated burst of length b . If the first non-zero component is the i th component, $i = 1, 2, \dots, n - 2b + 1$, then the first burst may go upto $(i + b - 1)$ th components and the second burst is in the remaining last $(n - i - b + 1)$ components. However, for $i = n - 2b + 2, \dots, n$ (i.e. in the last $2b - 1$ components) any number of non-zero components will be considered as a 2-repeated burst of length b or less.

The paper has been organised as follows. In Section 2, we derive a bound analogous to Fire's bound (see [6]; [11, Theorem 4.16]) for correction of 2-repeated burst of length b or less and then obtain in the following result another bound proving the existence of such a code. Section 3 describes bounds for m -repeated burst errors. Firstly, we obtain a necessary condition for the correction of m -repeated burst error correcting code followed by a sufficient condition for the existence of such a code. The last section, namely section 4 presents yet another bound for codes detecting m -repeated burst errors. The paper concludes with an illustration of a 2-repeated burst of length 3 or less correcting code and another illustration of a 2-repeated burst of length b or less correcting code over $\text{GF}(2)$.

In what follows, an (n, k) linear code will be considered as a subspace of the space of all n -tuples over $\text{GF}(q)$. The distance between two vectors shall be considered in the Hamming sense.

2. Correction of 2-Repeated Bursts

In this section, we consider linear codes capable of correcting 2-repeated bursts of length b or less. In the following, we obtain a necessary condition for the existence of such a code. The technique followed is that of the well-known Hamming's sphere-packing bound (see [7]; [11, Theorem 4.5]) of enumerating the total number of correctable error vectors and then comparing it with the total number of cosets.

Before we prove the main results, we give below an extension of problem 2.2(b) [9] a result which will be used in proving subsequent results of this paper. The problem stated in [9] is as follows:

Problem. Show that

$$\sum_{k=0}^{n-1} \binom{2+k}{2} = \binom{2+n}{3}$$

and thus

$$1 \times 2 + 2 \times 3 + 3 \times 4 + \dots + n(n+1) = \frac{1}{3}n(n+1)(n+2).$$

This result can be extended as follows:

$$\sum_{i=r}^n \binom{i}{r} = \binom{r}{r} + \binom{r+1}{r} + \dots + \binom{n}{r} = \binom{n+1}{r+1}. \quad (2.1)$$

Proof. We prove the result by induction on n .

For $n = 1$, r can only be 1.

L.H.S. = R.H.S. = 1.

Therefore the result holds for $n = 1$.

Assume the result to be true for $n = p$.

For $n = p + 1$,

$$\begin{aligned} & \binom{r}{r} + \binom{r+1}{r} + \dots + \binom{p+1}{r} \\ &= \left(\binom{r}{r} + \binom{r+1}{r} + \dots + \binom{p}{r} \right) + \binom{p+1}{r} \\ &= \binom{p+1}{r+1} + \binom{p+1}{r} \text{ (by induction hypothesis)} \\ &= \binom{p+2}{r+1} \text{ (using the identity } \binom{n}{r} + \binom{n}{r-1} = \binom{n+1}{r} \text{)}. \end{aligned}$$

Therefore the result holds for $n = p + 1$.

Hence, by the principle of mathematical induction, (2.1) is true for every positive integer n . $\square$

Alternatively, the above result may also be derived as follows:

Alternate Proof. Consider the finite geometric series

$$(1+x) + (1+x)^2 + (1+x)^3 + \dots + (1+x)^n. \quad (2.2)$$

The coefficient of x^r in (2.2) is

$$\binom{r}{r} + \binom{r+1}{r} + \binom{r+2}{r} + \dots + \binom{n}{r}. \quad (2.3)$$

Also the sum of the series (2.2) is

$$(1+x) \left(\frac{(1+x)^{n+1} - 1}{(1+x) - 1} \right) = \frac{(1+x)^{n+1} - (1+x)}{x}$$

which gives the coefficient of x^r as

$$\binom{n+1}{r+1}. \quad (2.4)$$

From (2.3) and (2.4), we get

$$\sum_{i=r}^n \binom{i}{r} = \binom{r}{r} + \binom{r+1}{r} + \binom{r+2}{r} + \dots + \binom{n}{r} = \binom{n+1}{r+1}. \quad \square$$

Theorem 2.1. Any (n, k) linear code over $\text{GF}(q)$ that corrects all 2-repeated bursts of length b or less must satisfy

$$q^{n-k} \geq q^{2(b-1)} \left(\binom{n-2b+2}{2} (q-1)^2 + \binom{n-2b+1}{1} (q-1) + q \right). \quad (2.5)$$

Proof. We first enumerate the total number of correctable error vectors which are 2-repeated bursts of length b or less.

Consider a 2-repeated burst of length b or less. Let the first non-zero component of the vector be the i th component which may go upto the $(i+b-1)$ th component, where $i = 1, 2, \dots, n-2b+1$. The second vector of length b or less will then be in the last (remaining) $(n-b-i+1)$ components. From Fire's bound (see [6]; [11, Theorem 4.16]), the number of bursts of length b or less in a vector of length $(n-b-i+1)$ is

$$\begin{aligned} & q^{b-1} \left(\binom{n-b-i+1-b+1}{1} (q-1) + 1 \right) \\ &= q^{b-1} \left(\binom{n-2b-i+2}{1} (q-1) + 1 \right). \end{aligned}$$

Summing over i gives the number of 2-repeated burst vectors with the first non-zero component in the first $(n-2b+1)$ components as

$$\begin{aligned} & \sum_{i=1}^{n-2b+1} (q-1) q^{b-1} \left[q^{b-1} \left(\binom{n-2b-i+2}{1} (q-1) + 1 \right) \right] \\ &= q^{2(b-1)} \left(\sum_{i=1}^{n-2b+1} \binom{n-2b-i+2}{1} (q-1)^2 + \sum_{i=1}^{n-2b+1} (q-1) \right) \\ &= q^{2(b-1)} \left(\binom{n-2b+2}{2} (q-1)^2 + \binom{n-2b+1}{1} (q-1) \right) \quad (\text{using (2.1)}). \end{aligned}$$

Also, the number of 2-repeated bursts of length b or less in the last $2b-1$ components is clearly q^{2b-1} . Thus, the total number of 2-repeated bursts of length b or less, including the vector of all zeros, is

$$\begin{aligned} & q^{2(b-1)} \left(\binom{n-2b+2}{2} (q-1)^2 + \binom{n-2b+1}{1} (q-1) \right) + q^{2b-1} \\ &= q^{2(b-1)} \left(\binom{n-2b+2}{2} (q-1)^2 + \binom{n-2b+1}{1} (q-1) + q \right). \end{aligned}$$

Since all such vectors being correctable error vectors must belong to distinct cosets of the standard array and the total number of cosets available equals q^{n-k} , we must have

$$q^{n-k} \geq q^{2(b-1)} \left(\binom{n-2b+2}{2} (q-1)^2 + \binom{n-2b+1}{1} (q-1) + q \right). \quad \square$$

In the following result, we derive a sufficient condition on the number of check digits required for the existence of such a code. The proof is based on the technique used to establish Varsharmov-Gilbert-Sacks bound by constructing a parity-check matrix for the desired code (see [11, Theorem 4.17]). This technique not only ensures the existence of such a code but also gives a method for constructing such codes.

Theorem 2.2. *There shall always exist an (n, k) linear code over $\text{GF}(q)$ that corrects all 2-repeated bursts of length b or less ($n > 4b$) provided that*

$$q^{n-k} > q^{4(b-1)} \left(\binom{n-4b+3}{3} (q-1)^3 + \binom{n-4b+2}{2} (q-1)^2 + \binom{n-4b+1}{1} (q-1)q + q^2 \right). \quad (2.6)$$

Proof. In order to prove this result, we first prove the following lemma.

Lemma 2.1. *The number of 3-repeated bursts of length b or less in a vector of length n is*

$$q^{3(b-1)} \left(\binom{n-3b+3}{3} (q-1)^3 + \binom{n-3b+2}{2} (q-1)^2 + \binom{n-3b+1}{1} (q-1)q + q^2 \right). \quad (2.7)$$

Proof. Consider a 3-repeated burst of length b or less. Let the first non-zero component of the vector be the i th component and such a burst may go upto the $(i+b-1)$ th component, where $i = 1, 2, \dots, n-3b+1$. The second and the third bursts of length b or less will then be in the last (remaining) $(n-b-i+1)$ components. To enumerate the number of second and third bursts of length b or less is equivalent to enumerating the number of 2-repeated bursts of length b or less in a vector of length $(n-b-i+1)$, which from Theorem 2.1 is

$$\begin{aligned} & q^{2(b-1)} \left(\binom{n-b-i+1-2b+2}{2} (q-1)^2 + \binom{n-b-i+1-2b+1}{1} (q-1) + q \right) \\ &= q^{2(b-1)} \left(\binom{n-3b-i+3}{2} (q-1)^2 + \binom{n-3b-i+2}{1} (q-1) + q \right). \end{aligned}$$

Summing on i gives the number of 3-repeated bursts of length b or less with the

first non-zero component in the first $n - 3b + 1$ components as

$$\begin{aligned}
& \sum_{i=1}^{n-3b+1} (q-1)q^{b-1} \left[q^{2(b-1)} \binom{n-3b-i+3}{2} (q-1)^2 \right. \\
& \quad \left. + \binom{n-3b-i+2}{1} (q-1) + q \right] \\
&= q^{3(b-1)} \left(\sum_{i=1}^{n-3b+1} \binom{n-3b+3-i}{2} (q-1)^3 \right. \\
& \quad \left. + \sum_{i=1}^{n-3b+1} \binom{n-3b+2-i}{1} (q-1)^2 + \sum_{i=1}^{n-3b+1} (q-1)q \right) \\
&= q^{3(b-1)} \left(\binom{n-3b+3}{3} (q-1)^3 + \binom{n-3b+2}{2} (q-1)^2 \right. \\
& \quad \left. + \binom{n-3b+1}{1} (q-1)q \right) \quad (\text{using (2.1)}).
\end{aligned}$$

Also, the number of 3-repeated bursts of length b or less in the last $3b-1$ components is clearly q^{3b-1} .

Thus the total number of 3-repeated bursts of length b or less in a vector of length n , including the vector of all zeros, is

$$\begin{aligned}
& q^{3(b-1)} \left(\binom{n-3b+3}{3} (q-1)^3 + \binom{n-3b+2}{2} (q-1)^2 \right. \\
& \quad \left. + \binom{n-3b+1}{1} (q-1)q \right) + q^{3b-1} \\
&= q^{3(b-1)} \left(\binom{n-3b+3}{3} (q-1)^3 + \binom{n-3b+2}{2} (q-1)^2 \right. \\
& \quad \left. + \binom{n-3b+1}{1} (q-1)q + q^2 \right)
\end{aligned}$$

which proves the lemma. $\square$

Proof of the theorem. We shall prove the result by constructing an appropriate $(n-k) \times n$ parity-check matrix $H = [h_1, h_2, \dots, h_n]$ for the desired code. Choose any non-zero $(n-k)$ -tuple as the first column h_1 of H . Suppose that we have selected the first $j-1$ columns $h_1, h_2, \dots, h_{j-1}$ of H suitably. We lay down the condition to add the j th column h_j as follows:

h_j should not be a linear combination of the immediately preceding $(b-1)$ or fewer columns $h_{j-b+1}, \dots, h_{j-1}$ of H together with any three distinct sets of b or fewer consecutive columns each from amongst the first $(j-b)$ columns $h_1, h_2, \dots, h_{j-b}$.

In other words,

$$\begin{aligned} h_j \neq & (\alpha_1 h_{j-b+1} + \alpha_2 h_{j-b+2} + \dots + \alpha_{b-1} h_{j-1}) \\ & + (\beta_1 h_{i_1} + \beta_2 h_{i_1+1} + \dots + \beta_b h_{i_1+b-1}) \\ & + (\gamma_1 h_{i_2} + \gamma_2 h_{i_2+1} + \dots + \gamma_b h_{i_2+b-1}) \\ & + (\delta_1 h_{i_3} + \delta_2 h_{i_3+1} + \dots + \delta_b h_{i_3+b-1}) \end{aligned}$$

where

$$\alpha_i, \beta_i, \gamma_i, \delta_i \in \text{GF}(q), \quad i_1 + i_2 + i_3 + 3b - 3 \leq j - b. \quad (2.8)$$

This condition ensures that there would not be a code-word which is expressible as a sum or difference of two vectors each of which is a 2-repeated burst of length b or less.

We now enumerate all possible linear combinations on the R.H.S. of (2.8).

Clearly, the coefficients α_i 's which are $(b-1)$ in number can be chosen in $q^{(b-1)}$ ways.

To compute the coefficients β_i 's, γ_i 's and δ_i 's is equivalent to find the number of 3-repeated bursts of length b or less in a vector of length $j-b$, which by Lemma 2.1 is

$$\begin{aligned} q^{3(b-1)} & \left(\binom{j-4b+3}{3} (q-1)^3 + \binom{j-4b+2}{2} (q-1)^2 \right. \\ & \left. + \binom{j-4b+1}{1} (q-1)q + q^2 \right). \end{aligned}$$

Thus, the total number of linear combinations on the R.H.S. of (2.8) equals

$$\begin{aligned} & q^{b-1} \left[q^{3(b-1)} \left(\binom{j-4b+3}{3} (q-1)^3 + \binom{j-4b+2}{2} (q-1)^2 \right. \right. \\ & \left. \left. + \binom{j-4b+1}{1} (q-1)q + q^2 \right) \right] \\ & = q^{4(b-1)} \left(\binom{j-4b+3}{3} (q-1)^3 + \binom{j-4b+2}{2} (q-1)^2 \right. \\ & \left. + \binom{j-4b+1}{1} (q-1)q + q^2 \right). \quad (2.9) \end{aligned}$$

Since these many linear combinations cannot be put to be equal to h_j and at worst all the linear combinations computed in (2.9) might yield a distinct sum, therefore in view of the fact that the total number of $(n-k)$ -tuples is q^{n-k} , the j th column h_j can be added to H provided that $q^{n-k} > (2.9)$.

To obtain a code of length n , replace j by n in the above inequality and we get

$$\begin{aligned} q^{n-k} & > q^{4(b-1)} \left(\binom{n-4b+3}{3} (q-1)^3 + \binom{n-4b+2}{2} (q-1)^2 \right. \\ & \left. + \binom{n-4b+1}{1} (q-1)q + q^2 \right). \quad \square \end{aligned}$$

3. Correction of m -Repeated Burst Errors

In this section, we first consider linear codes capable of correcting m -repeated bursts of length b or less. In the following, we obtain a necessary condition for the existence of such a code. The technique followed is that of the well-known Hamming's sphere-packing bound (see [7]; [11, Theorem 4.5]) as used in Theorem 2.1.

Theorem 3.1. *Any (n, k) linear code over $\text{GF}(q)$ that corrects all m -repeated bursts of length b or less must satisfy*

$$q^{n-k} \geq q^{m(b-1)} \left(\binom{n-mb+m}{m} (q-1)^m + \sum_{l=0}^{m-1} \binom{n-mb+l}{l} (q-1)^l q^{m-1-l} \right). \quad (3.1)$$

Proof. We first enumerate the total number of correctable error vectors which are m -repeated bursts of length b or less.

We prove the result by induction on m applied to the R.H.S. of inequality (3.1).

For $m = 1$, expression on the R.H.S. of (3.1) becomes

$$q^{(b-1)} \left(\binom{n-b+1}{1} (q-1) + 1 \right)$$

which is the number of bursts of length b or less in a vector of length n as given in Fire's bound (see [6]; [11, Theorem 4.16]). So the result holds for $m = 1$.

Assume it to be true for $m = p$. Now we prove the result for $m = p+1$. Consider a $(p+1)$ -repeated burst of length b or less. Let the first non-zero component of the vector be the i th component which may go upto the $(i+b-1)$ th component, where $i = 1, 2, \dots, n - (p+1)b + 1$. The remaining p bursts of length b or less each will then be in the last (remaining) $(n - b - i + 1)$ components. To enumerate the number of such remaining p bursts is equivalent to enumerate the number of p -repeated bursts of length b or less in a vector of length $(n - b - i + 1)$. Since the result has been assumed to be true for $m = p$, so applying this to a vector of length $(n - b - i + 1)$ gives the required number as

$$\begin{aligned} & q^{p(b-1)} \left(\binom{n-b-i+1-pb+p}{p} (q-1)^p \right. \\ & \left. + \sum_{\ell=0}^{p-1} \binom{n-b-i+1-pb+\ell}{\ell} (q-1)^\ell q^{p-1-\ell} \right) \\ & = q^{p(b-1)} \left(\binom{n-(p+1)b-i+(p+1)}{p} (q-1)^p \right. \\ & \quad \left. + \sum_{\ell=0}^{p-1} \binom{n-(p+1)b-i+(\ell+1)}{\ell} (q-1)^\ell q^{p-1-\ell} \right). \end{aligned}$$

Summing on i gives the number of $(p+1)$ -repeated bursts with the first non-zero

component in the first $(n - (p + 1)b + 1)$ components as

$$\begin{aligned}
 & \sum_{i=1}^{n-(p+1)b+1} (q-1)q^{b-1} \left[q^{p(b-1)} \binom{n-(p+1)b-i+(p+1)}{p} (q-1)^p \right. \\
 & \left. + \sum_{l=0}^{p-1} \binom{n-(p+1)b-i+(l+1)}{l} (q-1)^l q^{p-1-l} \right] \\
 & = q^{(p+1)(b-1)} \left(\sum_{i=1}^{n-(p+1)b+1} \binom{n-(p+1)b+(p+1)-i}{p} (q-1)^{p+1} \right. \\
 & \left. + \sum_{i=1}^{n-(p+1)b+1} \sum_{l=0}^{p-1} \binom{n-(p+1)b+(l+1)-i}{l} (q-1)^{l+1} q^{p-1-l} \right). \quad (3.2)
 \end{aligned}$$

$$\begin{aligned}
 & \sum_{i=1}^{n-(p+1)b+1} \binom{n-(p+1)b+(p+1)-i}{p} (q-1)^{p+1} \\
 & = \binom{n-(p+1)b+(p+1)}{p+1} (q-1)^{p+1} \quad (\text{using (2.1)}). \\
 & \sum_{i=1}^{n-(p+1)b+1} \sum_{l=0}^{p-1} \binom{n-(p+1)b+(l+1)-i}{l} (q-1)^{l+1} q^{p-1-l} \\
 & = \sum_{l=0}^{p-1} \sum_{i=1}^{n-(p+1)b+1} \binom{n-(p+1)b+(l+1)-i}{l} (q-1)^{l+1} q^{p-1-l} \\
 & = \sum_{l=0}^{p-1} \binom{n-(p+1)b+(l+1)}{l+1} (q-1)^{l+1} q^{p-1-l} \quad (\text{using (2.1)}) \\
 & = \sum_{l=1}^p \binom{n-(p+1)b+l}{l} (q-1)^l q^{p-l} \quad (\text{replacing } l+1 \text{ by } l).
 \end{aligned}$$

Then (3.2) becomes

$$\begin{aligned}
 & q^{(p+1)(b-1)} \left(\binom{n-(p+1)b+(p+1)}{p+1} (q-1)^{p+1} \right. \\
 & \left. + \sum_{l=1}^p \binom{n-(p+1)b+l}{l} (q-1)^l q^{p-l} \right).
 \end{aligned}$$

Also, the number of $(p + 1)$ -repeated bursts in the last $(p + 1)b - 1$ components is clearly $q^{(p+1)b-1}$.

Thus the total number of $(p + 1)$ -repeated bursts of length b or less in a vector

of length n is

$$\begin{aligned}
 & q^{(p+1)(b-1)} \left(\binom{n - (p+1)b + (p+1)}{p+1} (q-1)^{p+1} \right. \\
 & \left. + \sum_{l=1}^p \binom{n - (p+1)b + l}{l} (q-1)^l q^{p-l} \right) + q^{(p+1)b-1} \\
 & = q^{(p+1)(b-1)} \left(\binom{n - (p+1)b + (p+1)}{p+1} (q-1)^{p+1} \right. \\
 & \left. + \sum_{l=0}^p \binom{n - (p+1)b + l}{l} (q-1)^l q^{p-l} \right).
 \end{aligned}$$

Thus the result holds for $m = (p+1)$.

Hence, by the principle of mathematical induction, the number of m -repeated bursts of length b or less in a vector of length n is

$$q^{m(b-1)} \left(\binom{n - mb + m}{m} (q-1)^m + \sum_{l=0}^{m-1} \binom{n - mb + l}{l} (q-1)^l q^{m-1-l} \right).$$

The rest of the result follows as in Theorem 2.1. □

Remark 3.1. For $m = 1$, the expression obtained in (3.1) reduces to

$$q^{n-k} \geq q^{b-1} \left(\binom{n-b+1}{1} (q-1) + 1 \right)$$

which is the well-known Fire's bound (see [6]; [11, Theorem 4.16]).

Remark 3.2. For $m = 2$, the expression obtained in (3.1) reduces to

$$q^{n-k} \geq q^{2(b-1)} \left(\binom{n-2b+2}{2} (q-1)^2 + \sum_{l=0}^1 \binom{n-2b+l}{l} (q-1)^l q^{1-l} \right)$$

which coincides with the result obtained in Theorem 2.1 for 2-repeated bursts.

In the following result we derive a sufficient condition on the number of check digits required for the existence of m -repeated burst error correcting code. The proof is based on the technique used to establish Varsharmov-Gilbert-Sacks bound by constructing a parity-check matrix for the desired code (see [11, Theorem 4.17]). This technique not only ensures the existence of such a code but also gives a method for constructing such codes.

Theorem 3.2. *There shall always exist an (n, k) linear code over $\text{GF}(q)$ that cor-*

rects all m -repeated bursts of length b or less ($n > 2mb$) provided that

$$q^{n-k} > q^{2m(b-1)} \left(\binom{n-2mb+(2m-1)}{2m-1} (q-1)^{2m-1} + \sum_{l=0}^{2m-2} \binom{n-2mb+l}{l} (q-1)^l q^{2m-2-l} \right). \quad (3.3)$$

Proof. We shall prove the result by constructing an appropriate $(n-k) \times n$ parity-check matrix $H = [h_1, h_2, \dots, h_n]$ for the desired code. Choose any non-zero $(n-k)$ -tuple as the first column h_1 of H . Suppose that we have selected the first $j-1$ columns $h_1, h_2, \dots, h_{j-1}$ of H suitably. We lay down the condition to add the j th column h_j as follows:

h_j should not be a linear combination of the immediately preceding $(b-1)$ or fewer columns $h_{j-b+1}, \dots, h_{j-1}$ of H together with any $(2m-1)$ distinct sets of b or fewer consecutive columns each from amongst the first $(j-b)$ columns $h_1, h_2, \dots, h_{j-b}$.

In other words,

$$\begin{aligned} h_j &\neq (\alpha_1 h_{j-b+1} + \alpha_2 h_{j-b+2} + \dots + \alpha_{b-1} h_{j-1}) \\ &\quad + (\beta_1 h_{i_1} + \beta_2 h_{i_1+1} + \dots + \beta_b h_{i_1+b-1}) \\ &\quad + (\gamma_1 h_{i_2} + \gamma_2 h_{i_2+1} + \dots + \gamma_b h_{i_2+b-1}) + \dots \\ &\quad + (\delta_1 h_{i_{2m-1}} + \delta_2 h_{i_{2m-1}+1} + \dots + \delta_b h_{i_{2m-1}+b-1}) \end{aligned} \quad (3.4)$$

where $\alpha_i, \beta_i, \gamma_i, \dots, \delta_i \in \text{GF}(q)$ and $i_1 + i_2 + \dots + i_{2m-1} + (2m-1)b - (2m-1) \leq j-b$. This condition ensures that there would not be a code-word which is expressible as a sum or difference of two vectors each of which is an m -repeated burst of length b or less.

We now enumerate all possible combinations on the R.H.S. of (3.4).

Clearly, the coefficients α_i 's which are $(b-1)$ in number can be chosen in q^{b-1} ways.

To compute the coefficients β_i 's, γ_i 's, $\dots$, δ_i 's is equivalent to finding the number of $(2m-1)$ -repeated bursts of length b or less in a vector of length $j-b$, which by Theorem 3.1 is

$$\begin{aligned} q^{(2m-1)(b-1)} &\left(\binom{j-2mb+2m-1}{2m-1} (q-1)^{2m-1} \right. \\ &\quad \left. + \sum_{l=0}^{2m-2} \binom{j-2mb+l}{l} (q-1)^l q^{2m-2-l} \right). \end{aligned}$$

Thus, the total number of linear combinations on the R.H.S. of (3.4) equals

$$\begin{aligned}
 & q^{b-1} \left[q^{(2m-1)(b-1)} \binom{j-2mb+2m-1}{2m-1} (q-1)^{2m-1} \right. \\
 & \left. + \sum_{l=0}^{2m-2} \binom{j-2mb+l}{l} (q-1)^l q^{2m-2-l} \right] \\
 & = q^{2m(b-1)} \left(\binom{j-2mb+2m-1}{2m-1} (q-1)^{2m-1} \right. \\
 & \quad \left. + \sum_{l=0}^{2m-2} \binom{j-2mb+l}{l} (q-1)^l q^{2m-2-l} \right). \tag{3.5}
 \end{aligned}$$

Since these many linear combinations cannot be put to be equal to h_j and at worst all the linear combinations computed in (3.5) might yield a distinct sum, therefore in view of the fact that total number of $(n-k)$ -tuples is q^{n-k} , the j th column h_j can be added to H provided that $q^{n-k} > (3.5)$.

To obtain a code of length n , replace j by n in the above inequality and we get

$$\begin{aligned}
 q^{n-k} & > q^{2m(b-1)} \left(\binom{n-2mb+(2m-1)}{2m-1} (q-1)^{2m-1} \right. \\
 & \quad \left. + \sum_{l=0}^{2m-2} \binom{n-2mb+l}{l} (q-1)^l q^{2m-2-l} \right). \quad \square
 \end{aligned}$$

Remark 3.3. For $m = 1$, the expression obtained in (3.3) reduces to

$$q^{n-k} > q^{2(b-1)} \left(\binom{n-2b+1}{1} (q-1) + 1 \right)$$

which is sufficient condition of correction of (usual) bursts of length b or less (see [11, Theorem 4.17]).

Remark 3.4. For $m = 2$, the expression obtained in (3.3) reduces to

$$\begin{aligned}
 q^{n-k} & > q^{4(b-1)} \left(\binom{n-4b+3}{3} (q-1)^3 + \binom{n-4b+2}{2} (q-1)^2 \right. \\
 & \quad \left. + \binom{n-4b+1}{1} (q-1)q + q^2 \right)
 \end{aligned}$$

which coincides with the result obtained in Theorem 2.2 for 2-repeated bursts.

4. A Sufficient Condition for the Detection of m -Repeated Bursts

In this section, we derive upper bound on the number of parity-check digits required for the detection of m -repeated burst error detecting codes. The technique followed is the same as in Theorem 2.2. This technique not only ensures the existence of such a code but also gives a method for constructing such a code.

Theorem 4.1. *There shall always exist an (n, k) linear code over $\text{GF}(q)$ that has no m -repeated burst of length b or less as a code word provided that*

$$q^{n-k} > q^{m(b-1)} \left(\binom{n-mb+(m-1)}{m-1} (q-1)^{m-1} + \sum_{l=0}^{m-2} \binom{n-mb+l}{l} (q-1)^l q^{m-2-l} \right). \quad (4.1)$$

Proof. We shall prove the result by constructing an appropriate $(n-k) \times n$ parity-check matrix $H = [h_1, h_2, \dots, h_n]$ for the desired code. Choose any non-zero $(n-k)$ -tuple as the first column h_1 of H . Suppose that we have selected the first $j-1$ columns $h_1, h_2, \dots, h_{j-1}$ of H suitably. We lay down the condition to add the j th column h_j as follows:

h_j should not be a linear combination of the immediately preceding $(b-1)$ or fewer columns $h_{j-b+1}, h_{j-b+2}, \dots, h_{j-1}$ of H together with any $(m-1)$ distinct sets of b or fewer consecutive columns each from amongst the first $j-b$ columns $h_1, h_2, \dots, h_{j-b}$.

In other words,

$$\begin{aligned} h_j &\neq (\alpha_1 h_{j-b+1} + \alpha_2 h_{j-b+2} + \dots + \alpha_{b-1} h_{j-1}) \\ &\quad + (\beta_1 h_{i_1} + \beta_2 h_{i_1+1} + \dots + \beta_b h_{i_1+b-1}) \\ &\quad + (\gamma_1 h_{i_2} + \gamma_2 h_{i_2+1} + \dots + \gamma_b h_{i_2+b-1}) + \dots \\ &\quad + (\delta_1 h_{i_{m-1}} + \delta_2 h_{i_{m-1}+1} + \dots + \delta_b h_{i_{m-1}+b-1}) \end{aligned}$$

where $\alpha_i, \beta_i, \gamma_i, \dots, \delta_i \in \text{GF}(q)$ and

$$i_1 + i_2 + \dots + i_{m-1} + (m-1)b - (m-1) \leq j-b. \quad (4.2)$$

This condition ensures that no m -repeated burst of length b or less will be a code word.

We now enumerate all possible linear combinations on the R.H.S. of (4.2).

Clearly, the coefficients α_i 's which are $(b-1)$ in number can be chosen in q^{b-1} ways.

To compute the coefficients β_i 's, γ_i 's, $\dots$, δ_i 's is equivalent to find the number of $(m-1)$ -repeated bursts of length b or less in a vector of length $j-b$, which by Theorem 3.1 is

$$\begin{aligned} q^{(m-1)(b-1)} &\left(\binom{j-mb+(m-1)}{m-1} (q-1)^{m-1} \right. \\ &\quad \left. + \sum_{l=0}^{m-2} \binom{j-mb+l}{l} (q-1)^l q^{m-2-l} \right). \end{aligned}$$

Thus, the total number of linear combinations on the R.H.S. of (4.2) equals

$$\begin{aligned}
 & q^{b-1} \left[q^{(m-1)(b-1)} \binom{j-mb+(m-1)}{(m-1)} (q-1)^{m-1} \right. \\
 & \quad \left. + \sum_{l=0}^{m-2} \binom{j-mb+l}{l} (q-1)^l q^{m-2-l} \right] \\
 & = q^{m(b-1)} \left(\binom{j-mb+(m-1)}{(m-1)} (q-1)^{m-1} \right. \\
 & \quad \left. + \sum_{l=0}^{m-2} \binom{j-mb+l}{l} (q-1)^l q^{m-2-l} \right). \tag{4.3}
 \end{aligned}$$

Since these many linear combinations cannot be put to be equal to h_j and at worst all the linear combinations computed in (4.3) might yield a distinct sum, therefore in view of the fact that the total number of $(n-k)$ -tuples is q^{n-k} , the j th column h_j can be added to H provided that $q^{n-k} > (4.3)$.

To obtain a code of length n , replace j by n in the above inequality and we get

$$\begin{aligned}
 q^{n-k} & > q^{m(b-1)} \left(\binom{n-mb+(m-1)}{(m-1)} (q-1)^{m-1} \right. \\
 & \quad \left. + \sum_{l=0}^{m-2} \binom{n-mb+l}{l} (q-1)^l q^{m-2-l} \right). \quad \square
 \end{aligned}$$

Remark 4.1. For $m = 1$ the expression obtained in (4.1) reduces to

$$q^{n-k} > q^{b-1}$$

which is the sufficient condition of detection of (usual) bursts of length b or less (see [11, Theorem 4.14]).

Remark 4.2. For $m = 2$, the expression obtained in (4.1) reduces to

$$q^{n-k} > q^{2(b-1)} \left(\binom{n-2b+1}{1} (q-1) + 1 \right)$$

which coincides with the result obtained in [2, Theorem 2] for 2-repeated bursts.

Remark 4.3. We note that the necessary condition for correction of m -repeated bursts (Theorem 3.3) is required to be proved first because it is needed to prove the sufficient condition of detection of m -repeated bursts (Theorem 4.1).

Remark 4.4. In view of the fact that the result obtained in Theorem 4.1 for the detection of $2m$ -repeated bursts is the same as the result for the correction of m -repeated bursts of length b or less, it is worthwhile to note that such a code can serve dual purpose, i.e., it can either be used to correct m -repeated bursts of length b or less or it can be used to detect a $2m$ -repeated burst of length b or less.

We conclude with the following two examples.

Example 4.1. Consider a $(15, 3)$ binary code with parity-check matrix

$$H = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 \end{bmatrix}_{12 \times 15}$$

This matrix has been constructed by the synthesis procedure outlined in the proof of Theorem 2.2 by taking $b = 3$ over $\text{GF}(2)$. It can be seen from Table 1 that the syndromes of different 2-repeated bursts of length 3 or less are distinct, showing thereby that the code that is the null space of the matrix given in the example, corrects all 2-repeated bursts of length 3 or less.

Table 1. Error Pattern – Syndrome Table.

Error vectors	Syndromes	Error vectors	Syndromes
1000000000000000	10000000000000	1000011100000000	10000111000000
1001000000000000	10010000000000	1000001110000000	10000011100000
1000100000000000	10001000000000	1000000111000000	10000001110000
1000010000000000	10000100000000	1000000011100000	10000000111000
1000001000000000	10000010000000	1000000001110000	10000000011100
1000000100000000	10000001000000	1000000000111000	10000000001110
1000000010000000	10000000100000	1000000000011100	00010010011111
1000000001000000	10000000010000	1000000000001110	01011011011011
1000000000100000	10000000001000	1000000000000111	01111111111111
1000000000010000	10000000000100	1100000000000000	11000000000000
1000000000001000	10000000000010	1101000000000000	11010000000000
1000000000000100	00010010010010	1100100000000000	11001000000000
1000000000000010	11001001001010	1100010000000000	11000100000000
1000000000000001	10100100100100	1100001000000000	11000010000000
1001100000000000	10011000000000	1100000100000000	11000001000000
1000110000000000	10001100000000	1100000010000000	11000000100000
1000011000000000	10000110000000	1100000001000000	11000000010000
1000001100000000	10000011000000	1100000000100000	11000000001000
1000000110000000	10000001100000	1100000000010000	01010010010010
1000000011000000	10000000110000	1100000000000100	10001001001001
1000000001100000	10000000011000	1100000000000001	11100100100100
1000000000110000	00010010010010	1101100000000000	11011000000000
1000000000011000	01011011011010	1100110000000000	11001100000000
1000000000001100	11101101101101	1100011000000000	11000110000000
1001010000000000	10010100000000	1100001100000000	11000011000000
1000101000000000	10001010000000	1100000110000000	11000001100000
1000010100000000	10000101000000	1100000011000000	11000000110000
1000001010000000	10000010100000	1100000001100000	11000000011000
1000000101000000	10000001010000	1100000000110000	01010010010010
1000000010100000	10000000101000	1100000000011000	00011011011010
1000000001010000	00010010010010	1100000000000110	10101101101101
1000000000101000	11001001001001	1101010000000000	11010100000000
1000000000010100	00110110110101	1100101000000000	11001010000000
1001110000000000	10011100000000	1100010100000000	11000101000000
1000111000000000	10001110000000	1100001010000000	11000010100000

Table 1 (Continued)

Error vectors	Syndromes
1100000010100000	11000000101000
1100000001010000	11000000010100
1100000000101000	11000000001010
1100000000010100	01010001001110
1100000000001010	10001001001111
1100000000000101	01110110110111
1101110000000000	11011100000000
1100111000000000	11001110000000
1100011100000000	11000111000000
1100001110000000	11000011100000
1100000111000000	11000001110000
1100000011100000	11000000111000
1100000001110000	11000000011100
1100000000111000	01010010011111
1100000000011100	00011011011111
1100000000001110	00111111111111
1010000000000000	10100000000000
1011000000000000	10110000000000
1010100000000000	10101000000000
1010010000000000	10100100000000
1010001000000000	10100010000000
1010000100000000	10100001000000
1010000010000000	10100000100000
1010000001000000	10100000010000
1010000000100000	10100000001000
1010000000010000	10100000000100
1010000000001000	10100000000010
1010000000000100	00110010010010
1010000000000010	11101001001010
1010000000000001	10000100100101
1011100000000000	10111000000000
1010110000000000	10101100000000
1010011000000000	10100110000000
1010001100000000	10100011000000
1010000110000000	10100001100000
1010000011000000	10100000110000
1010000001100000	10100000011000
1010000000110000	10100000001100
1010000000011000	10100000000110
1010000000001100	00110010010010
1010000000000110	01111011011011
1010000000000011	01111011011011

Error vectors	Syndromes
1010000000000011	11001101101011
1011010000000000	10110100000000
1010101000000000	10101010000000
1010010100000000	10100101000000
1010001010000000	10100010100000
1010000101000000	10100001010000
1010000010100000	10100000101000
1010000001010000	10100000010100
1010000000101000	10100000001010
1010000000010100	00110010011010
1010000000001010	11101001001011
1010000000000101	00010110110101
1011110000000000	10111100000000
1010111000000000	10101110000000
1010011100000000	10100111000000
1010001110000000	10100011100000
1010000111000000	10100001110000
1010000011100000	10100000111000
1010000001110000	10100000011100
1010000000111000	00110010011111
1010000000011100	01111011011011
1010000000001110	01011111111111
1110000000000000	11100000000000
1111000000000000	11110000000000
1110100000000000	11101000000000
1110010000000000	11100100000000
1110001000000000	11100010000000
1110000100000000	11100001000000
1110000010000000	11100000100000
1110000001000000	11100000010000
1110000000100000	11100000001000
1110000000010000	01110010010010
1110000000001000	10101001001010
1110000000000100	11000100100101
1111100000000000	11111000000000
1110110000000000	11101100000000
1110011000000000	11100110000000
1110001100000000	11100011000000
1110000110000000	11100001100000
1110000011000000	11100000110000

Table 1 (*Continued*)

Error vectors	Syndromes	Error vectors	Syndromes
1110000001100000	1110000001100	010001100000000	010001100000
1110000000110000	1110000000110	010000110000000	010000110000
1110000000011000	1110000000011	010000011000000	010000011000
1110000000001100	011100100101	010000001100000	010000001100
1110000000000110	001110110110	010000000110000	010000000110
1110000000000011	100011011011	010000000011000	010000000011
1111010000000000	111101000000	010000000001100	110100100101
1110101000000000	111010100000	010000000000110	100110110110
1110010100000000	111001010000	010000000000011	001011011011
1110001010000000	111000101000	010010100000000	010010100000
1110000101000000	111000010100	010001010000000	010001010000
1110000010100000	111000001010	010000101000000	010000101000
1110000001010000	111000000101	010000010100000	010000010100
1110000000101000	011100100110	010000001010000	010000001010
1110000000010100	101010010011	010000000101000	010000000101
1110000000001010	010101101101	010000000010100	110100100110
1111110000000000	111111000000	010000000001010	000010010011
1110111000000000	111011100000	010000000000101	111101101101
1110011100000000	111001110000	010011100000000	010011100000
1110001110000000	111000111000	010001110000000	010001110000
1110000111000000	111000011100	010000111000000	010000111000
1110000011100000	111000001110	010000011100000	010000011100
1110000001110000	111000000111	010000001110000	010000001110
1110000000111000	011100100111	010000000111000	110100100111
1110000000011100	000111111111	010000000011100	100110110111
0100000000000000	010000000000	010000000000111	101111111111
0100100000000000	010010000000	011000000000000	011000000000
0100010000000000	010001000000	011010000000000	011010000000
0100001000000000	010000100000	011001000000000	011001000000
0100000100000000	010000010000	011000100000000	011000100000
0100000010000000	010000001000	011000010000000	011000010000
0100000001000000	010000000100	011000001000000	011000001000
0100000000100000	010000000010	011000000100000	011000000100
0100000000010000	010000000001	011000000010000	011000000010
0100000000001000	110100100100	011000000001000	011000000001
0100000000000100	000010010010	011000000000100	111100100100
0100000000000010	011001001001	011000000000010	001010010010
0100110000000000	010011000000	011000000000001	010001001001

Table 1 (*Continued*)

Error vectors	Syndromes
0110110000000000	01101100000000
0110011000000000	01100110000000
0110001100000000	01100011000000
0110000110000000	01100001100000
0110000011000000	01100000110000
0110000001100000	01100000011000
0110000000110000	01100000001100
0110000000011000	01100000000110
0110000000001100	11110010010101
0110000000000110	10111011011010
0110000000000011	00001101101101
0110101000000000	01101010000000
0110010100000000	01100101000000
0110001010000000	01100010100000
0110000101000000	01100001010000
0110000010100000	01100000101000
0110000001010000	01100000010100
0110000000101000	11110010011010
0110000000010100	00101001001011
0110000000001010	11010110110101
0110111000000000	01101110000000
0110011100000000	01100111000000
0110001110000000	01100011100000
0110000111000000	01100001110000
0110000011100000	01100000111000
0110000001110000	01100000011100
0110000000111000	11110010011011
0110000000011100	10111011011011
0110000000001110	10011111111111
0101000000000000	01010000000000
0101100000000000	01011000000000
0101010000000000	01010100000000
0101001000000000	01010010000000
0101000100000000	01010001000000
0101000010000000	01010000100000
0101000001000000	01010000010000
0101000000100000	01010000001000
0101000000010000	01010000000100
0101000000001000	01010000000010
0101000000000100	11000010010010
0101000000000010	00011001001010

Error vectors	Syndromes
0101000000000001	01110100100101
0101110000000000	01011100000000
0101011000000000	01010110000000
0101001100000000	01010011000000
0101000110000000	01010001100000
0101000011000000	01010000110000
0101000001100000	01010000011000
0101000000110000	01010000001100
0101000000011000	01010000000110
0101000000001100	11000010010101
0101000000000110	10001011011010
0101000000000011	00111101101101
0101101000000000	01011010000000
0101010100000000	01010101000000
0101001010000000	01010010100000
0101000101000000	01010001010000
0101000010100000	01010000101000
0101000001010000	01010000010100
0101000000101000	01010000001010
0101000000010100	11000010011010
0101000000001010	00011001001011
0101000000000101	11100110110101
0101111000000000	01011110000000
0101011100000000	01010111000000
0101001110000000	01010011100000
0101000111000000	01010001110000
0101000011100000	01010000111000
0101000001110000	01010000011100
0101000000111000	01010000001110
0101000000011100	01010000000111
0101000000001110	11000010011011
0101000000000111	10101111111111
0111000000000000	01110000000000
0111100000000000	01111000000000
0111010000000000	01110100000000
0111001000000000	01110010000000
0111000100000000	01110001000000
0111000010000000	01110000100000
0111000001000000	01110000010000
0111000000100000	01110000001000
0111000000010000	01110000000100
0111000000001000	01110000000010
0111000000000100	01110000000001
0111000000000010	11100010010010

Table 1 (*Continued*)

Error vectors	Syndromes
0111000000000010	001110010010
0111000000000001	010101001001
0111110000000000	011111000000
0111011000000000	011101100000
0111001100000000	011100110000
0111000110000000	011100011000
0111000011000000	011100001100
0111000001100000	011100000110
0111000000110000	011100000011
011100000011000	111000100101
011100000001100	101010110110
011100000000110	000111011011
0111101000000000	011110100000
0111010100000000	011101010000
0111001010000000	011100101000
0111000101000000	011100010100
0111000010100000	011100001010
0111000001010000	011100000101
011100000010100	111000100110
011100000001010	001110010011
011100000000101	110001101101
0111111000000000	011111100000
0111011100000000	011101110000
0111001110000000	011100111000
0111000111000000	011100011100
0111000011100000	011100001110
011100000111000	011100000111
011100000011100	111000100111
011100000001110	101010110111
011100000000111	100011111111
0010000000000000	001000000000
0010010000000000	001001000000
0010001000000000	001000100000
0010000100000000	001000010000
0010000010000000	001000001000
0010000001000000	001000000100
001000000010000	001000000010
001000000001000	001000000001
001000000000100	101000100100
001000000000010	0010000000010
001000000000001	000101001001
001000000000000	001101100000

Error vectors	Syndromes
0010000000000010	011010010010
0010000000000001	000001001001
0010011000000000	001001100000
0010001100000000	001000110000
0010000110000000	001000011000
0010000011000000	001000001100
0010000001100000	001000000110
001000000011000	001000000011
001000000001100	101100100101
001000000000110	111110110110
001000000000011	010011011011
0010010100000000	001001010000
0010001010000000	001000101000
0010000101000000	001000010100
0010000010100000	001000001010
001000000101000	001000000101
001000000010100	101100100110
001000000001010	011010010011
001000000000101	100101101101
0010011100000000	001001110000
0010001110000000	001000111000
0010000111000000	001000011100
0010000011100000	001000001110
001000000111000	001000000111
001000000011100	101100100111
001000000001110	111110110111
001000000000111	110111111111
0011000000000000	001100000000
0011010000000000	001101000000
0011001000000000	001100100000
0011000100000000	001100010000
0011000010000000	001100001000
0011000001000000	001100000100
001100000010000	001100000010
001100000001000	001100000001
001100000000100	101000100100
001100000000010	011110010010
001100000000001	000101001001
0011011000000000	001101100000

Table 1 (*Continued*)

Error vectors	Syndromes
001100110000000	0011001100000
001100011000000	0011000110000
001100001100000	0011000011000
001100000110000	0011000001100
001100000011000	0011000000110
001100000001100	0011000000011
0011000000001100	101000100101
0011000000000110	111010110110
0011000000000011	010111011011
001101010000000	0011010100000
001100101000000	0011001010000
001100010100000	0011000101000
001100010100000	0011000101000
001100001010000	0011000010100
001100001010000	0011000010100
001100000101000	0011000001010
001100000101000	101000100110
001100000010100	011110010011
001100000001010	100001101101
001101110000000	0011011100000
001100111000000	0011001110000
001100011100000	0011000111000
001100001110000	0011000011100
001100000111000	0011000001110
001100000011100	101000100111
001100000001110	111010110111
001100000000111	110011111111
001010000000000	0010100000000
001011000000000	0010110000000
001010100000000	0010101000000
001010010000000	0010100100000
001010001000000	0010100010000
001010000100000	0010100001000
001010000010000	0010100000100
001010000001000	0010100000010
001010000000100	0010100000001
001010000000010	101110100100
001010000000001	000011001001
001011100000000	0010111000000
001010110000000	0010101100000
001010011000000	0010100110000
001010001100000	0010100011000
001010000110000	0010100001100
001010000011000	0010100000110
001010000001100	101010100101

Error vectors	Syndromes
0010100000110000	0010100000110
0010100000011000	0010100000011
0010100000001100	101110100101
0010100000000110	111100110110
0010100000000011	010001011011
0010110100000000	0010110100000
00101010101000000	00101010101000
00101001010100000	00101001010100
00101000101010000	00101000101010
00101000010101000	00101000010101
00101000001010100	101110100110
00101000000101010	011000010011
00101000000010101	100111101101
0010111100000000	0010111100000
0010101110000000	0010101110000
0010100111000000	0010100111000
0010100011100000	0010100011100
0010100001110000	0010100001110
0010100000111000	0010100000111
0011100000000000	0011100000000
0011110000000000	0011110000000
0011101000000000	0011101000000
0011100100000000	0011100100000
0011100010000000	0011100010000
0011100001000000	0011100001000
0011100000100000	0011100000100
0011100000010000	0011100000010
0011100000001000	101010100100
0011100000000100	011100010010
0011100000000010	000111001001
0011111000000000	0011111000000
0011101100000000	0011101100000
0011100110000000	0011100110000
0011100011000000	0011100011000
0011100001100000	0011100001100
0011100000110000	0011100000110
0011100000011000	101010100101

Table 1 (*Continued*)

[illegible]

Error vectors	Syndromes
000110000111000	0001100000111
000110000011100	100010100111
000110000001110	110000110111
000110000000111	111001111111
000101000000000	000101000000
000101100000000	000101100000
000101010000000	000101010000
000101001000000	000101001000
000101000100000	000101000100
000101000010000	000101000010
000101000001000	000101000001
000101000000100	100001100100
000101000000010	010111010010
000101000000001	001100001001
000101110000000	000101110000
000101011000000	000101011000
000101001100000	000101001100
000101000110000	000101000110
000101000011000	000101000011
000101000001100	100001100101
000101000000110	110011110110
000101000000011	011110011011
000101101000000	000101101000
000101001100000	000101001100
000101000110000	000101000110
000101000011000	000101000011
000101000001100	100001100110
000101000000110	110011110110
000101000000011	011110011011
000101101000000	000101101000
000101010100000	000101010100
000101001010000	000101001100
000101000101000	000101000101
000101000010100	100011100111
000101000001010	010111010011
000101000000101	101000101101
000101101100000	000101101000
000101010101000	000101010100
000101001011000	000101001100
000101000101100	000101000101
000101000010110	100011100111
000101000001011	010111010011
000101000000101	101000101101
000101101100000	000101101000
000101010101000	000101010100
000101001011000	000101001100
000101000101100	000101000101
000101000010110	100011100111
000101000001011	010111010011
000101000000101	101000101101
000010000000000	000010000000
000010010000000	000010010000
000010001000000	000010001000
000010000100000	000010000100
000010000010000	000010000010
000010000001000	000010000001
000010000000100	100110100100
000010000000010	010000010010
000010000000001	001011001001
000010011000000	000010011000

Error vectors	Syndromes
000111001000000	000111001000
000111000100000	000111000100
000111000010000	000111000010
000111000001000	000111000001
000111000000100	100011100100
000111000000010	010101010010
000111000000001	001110001001
000111110000000	000111110000
000111011000000	000111011000
000111001100000	000111001100
000111000110000	000111000110
000111000011000	000111000011
000111000001100	100011100101
000111000000110	110001110110
000111000000011	011100011011
000111101000000	000111101000
000111010100000	000111010100
000111001010000	000111001100
000111000101000	000111000101
000111000010100	100011100110
000111000001010	010101010011
000111000000101	101010101101
000111111000000	000111111000
000111011100000	000111011100
000111001110000	000111001110
000111000111000	000111000111
000111000011100	100011100111
000111000001110	110001110111
000111000000111	111000111111
000010000000000	000010000000
000010010000000	000010010000
000010001000000	000010001000
000010000100000	000010000100
000010000010000	000010000010
000010000001000	000010000001
000010000000100	100110100100
000010000000010	010000010010
000010000000001	001011001001
000010011000000	000010011000

Table 1 (*Continued*)

Error vectors	Syndromes	Error vectors	Syndromes
000010001100000	000010001100	000011000000101	101110101101
000010000110000	000010000110	000011011100000	000011011100
000010000011000	000010000011	000011001110000	000011001110
000010000001100	100110100101	000011000111000	000011000111
000010000000110	110100110110	000011000011100	100111100111
000010000000011	011001011011	000011000001110	110101110111
000010010100000	000010010100	000011000000111	111100111111
000010001010000	000010001010	000010100000000	000010100000
000010000101000	000010000101	000010110000000	000010110000
000010000010100	100110100110	000010101000000	000010101000
000010000001010	010000010011	000010100100000	000010100100
000010000000101	101111101101	000010100010000	000010100010
000010011100000	000010011100	000010100001000	000010100001
000010001110000	000010001110	000010100000100	100110000100
000010000111000	000010000111	000010100000010	010000110010
000010000011100	100110100111	000010100000001	001011101001
000010000001110	110100110111	000010111000000	000010111000
000010000000111	111101111111	000010101100000	000010101100
000011000000000	000011000000	000010100110000	000010100110
000011010000000	000011010000	000010100011000	000010100011
000011001000000	000011001000	000010100001100	100110000101
000011000100000	000011000100	000010100000110	110100010110
000011000010000	000011000010	000010100000011	011001111011
000011000001000	000011000001	000010110100000	000010110100
000011000000100	100111100100	000010101010000	000010101010
000011000000010	010001010010	000010100101000	000010100101
000011000000001	001010001001	000010100010100	100110000110
000011011000000	000011011000	000010100001010	010000110011
000011001100000	000011001100	000010100000101	101111001101
000011000110000	000011000110	000010111100000	000010111100
000011000011000	000011000011	000010101110000	000010101110
000011000001100	100111100101	000010100111000	000010100111
000011000000110	110101110110	000010100011100	100110000111
000011000000011	011000011011	000010100001110	110100010111
000011010100000	000011010100	000010100000111	111101011111
000011001010000	000011001010	000011100000000	000011100000
000011000101000	000011000101	000011110000000	000011110000
000011000010100	100111100110	000011101000000	000011101000
000011000001010	010001010011	000011100100000	000011100100

Table 1 (*Continued*)

Error vectors	Syndromes	Error vectors	Syndromes
00000010100000010	0100110000010	000001110011100	1001010101011
00000010100000001	0010000011001	000001110000110	1101110000111
00000010111000000	0000001011100	000001110000011	111110001111
00000010101100000	0000001010110	000000100000000	000000100000
00000010100110000	0000001010011	000000100100000	000000100100
00000010100011000	100101110101	000000100010000	000000100010
00000010100001100	110111100110	000000100001000	000000100001
00000010100000011	011010001011	000000100000100	100100000100
00000010110100000	0000001011010	000000100000010	010010110010
00000010101010100	0000001010101	000000100000001	001001101001
00000010100101010	100101110110	000000100110000	000000100110
00000010100010101	010011000011	000000100011000	000000100011
0000001010000101	101100111101	000000100001100	100100000101
00000010111100000	0000001011110	000000100000110	110110010110
00000010101110000	0000001010111	000000100000011	011011111011
00000010100111000	100101110111	000000100101000	000000100101
00000010100011100	110111100111	000000100010100	100100000110
0000001010000111	111110101111	000000100001010	010010110011
00000011100000000	0000001110000	000000100000101	101101001101
00000011110000000	0000001111000	000000100111000	000000100111
00000011101000000	0000001110100	000000100011100	100100000111
00000011100100000	0000001110010	000000100001110	110110010111
00000011100010000	0000001110001	000000100000111	111111011111
00000011100001000	1001010101010	000000110000000	000000110000
00000011100000010	010011100010	000000110100000	000000110100
00000011100000001	001000111001	000000110010000	000000110010
00000011111000000	0000001111100	000000110001000	000000110001
00000011101100000	0000001110110	000000110000100	100100010100
00000011100110000	0000001110011	000000110000010	010010100010
00000011100011000	100101010101	000000110000001	001001111001
00000011100001100	110111000110	000000110110000	000000110110
00000011100000011	011010101011	000000110011000	000000110011
00000011110100000	0000001111010	000000110001100	100100010101
00000011101010000	0000001110101	000000110000110	110110000110
00000011100101000	100101010110	000000110000011	011011101011
00000011100010100	010011100011	000000110101000	000000110101
0000001110000101	101100011101	000000110010100	100100010110
00000011111100000	0000001111110	000000110001010	010010100011
00000011101110000	0000001110111	000000110000101	101101011101

Error vectors	Syndromes	Error vectors	Syndromes
0000000010101000	000000001010101	000000001000101	101101100101
0000000010100100	100100110000	000000001001110	110110111111
0000000010100010	010010000110	000000001000111	111111110111
0000000010100001	0010010111101	000000001100000	000000001100
0000000010111000	0000000010111	000000001101000	000000001101
0000000010101100	100100110001	000000001100100	100100101000
0000000010100110	110110100010	000000001100010	010010011110
0000000010100011	011011001111	000000001100001	001001000101
0000000010110100	100100110010	000000001101100	100100101001
0000000010101010	010010000111	000000001100110	110110111010
0000000010100101	101101111001	000000001100011	011011010111
0000000010111100	100100110011	000000001101010	010010011111
0000000010101110	110110100011	000000001100101	101101100001
0000000010100111	111111101011	000000001101110	110110111011
0000000011100000	0000000011100	000000001100111	111111110011
0000000011110000	0000000011110	000000001010000	000000001010
0000000011101000	0000000011101	000000001011000	000000001011
0000000011100100	100100111000	000000001010100	100100101110
0000000011100010	010010001110	000000001010010	010010011000
0000000011100001	001001010101	000000001010001	001001000011
0000000011111000	0000000011111	000000001011100	100100101111
0000000011101100	100100111001	000000001010110	110110111100
0000000011100110	110110101010	000000001010011	011011010001
0000000011100011	011011000111	000000001011010	010010011001
0000000011110100	100100111010	000000001010101	101101100111
0000000011101010	010010001111	000000001011110	110110111101
0000000011100101	101101110001	000000001010111	111111110101
0000000011111100	100100111011	000000001110000	000000001110
0000000011101110	110110101011	000000001111000	000000001111
0000000011100111	111111100011	000000001110100	100100101010
0000000010000000	000000001000	000000001110010	010010011100
000000001001000	000000001001	000000001110001	001001000111
000000001000100	100100101100	000000001111100	100100101011
000000001000010	010010011010	000000001110110	110110111000
000000001000001	001001000001	000000001110011	011011010101
000000001001100	100100101101	000000001111010	010010011101
000000001000110	110110111110	000000001110101	101101100011
000000001000011	011011010011	000000001111110	110110111001
000000001001010	010010011011	000000001110111	111111110001

Table 1 (*Continued*)

Error vectors	Syndromes
0000000000100000	00000000000100
0000000000100100	1001001000000
0000000000100010	010010010110
0000000000100001	0010010011101
0000000000100110	110110110010
0000000000100011	011011011111
0000000000100101	101101101001
0000000000100111	111111111011
0000000000110000	0000000000110
0000000000110100	100100100010
0000000000110010	010010010100
0000000000110001	001001001111
0000000000110110	110110110000
0000000000110011	011011011101
0000000000110101	101101101011
0000000000110111	111111111001
0000000000101000	0000000000101
0000000000101100	100100100001
0000000000101010	010010010111
0000000000101001	001001001100
0000000000101110	110110110011
0000000000101011	011011011110
0000000000101101	101101101000
0000000000101111	111111111010
0000000000111000	0000000000111
0000000000111100	100100100011
0000000000111010	010010010101
0000000000111001	001001001110
0000000000111110	110110110001
0000000000111011	011011011100
0000000000111101	101101101010
0000000000111111	111111111000
0000000000011111	111111111100
0000000000011110	110110110101
0000000000011101	101101101110
0000000000011100	100100100111
0000000000011011	011011011000
0000000000011010	010010010001
0000000000011001	001001001010

Error vectors	Syndromes
0000000000011000	0000000000011
0000000000010111	111111111101
0000000000010110	110110110100
0000000000010101	101101101111
0000000000010100	100100100110
0000000000010011	011011011001
0000000000010010	010010010000
0000000000010001	001001001011
0000000000010000	0000000000010
0000000000001111	111111111110
0000000000001110	110110110111
0000000000001101	101101101100
0000000000001100	100100100101
0000000000001011	011011011010
0000000000001010	010010010011
0000000000001001	001001001000
0000000000001000	0000000000001
0000000000000111	111111111111
0000000000000110	110110110110
0000000000000101	101101101101
0000000000000100	100100100100
0000000000000011	011011011011
0000000000000010	010010010010
0000000000000001	001001001001

it is a linear combination of three distinct sets of b or less columns of H . Also, considering the first and the last b consecutive components, there will be $2^{b^2} = 2^{2b}$ patterns of components. The condition requires that any vector having first b or last b components same as a linear combination of $C_{4b+2}, C_{4b+3}, \dots, C_{5b}$ cannot be added to H .

Thus the vectors that could have been added to H are the ones which involve C_{4b+1} . On considering the linear combination of $C_{3b+2}, C_{3b+3}, \dots, C_{4b}, C_{4b+1}$ and $C_{4b+2}, C_{4b+3}, \dots, C_{5b}$, we find that all the remaining 2^{4b} tuples have their last b components as linear combination of $C_{3b+2}, \dots, C_{4b+1}, C_{4b+2}, \dots, C_{5b}$ and first b components as linear combination of $C_{4b+2}, \dots, C_{5b}$. In other words, the tuple has first b components same as a linear combination of $C_{4b+2}, \dots, C_{5b}$ and the last b components as linear combination of $C_{3b+2}, \dots, C_{4b+1}, C_{4b+2}, \dots, C_{5b}$ which gives us a linear combination of b consecutive columns $C_{3b+2}, \dots, C_{4b+1}$ and immediately preceding $(b-1)$ columns $C_{4b+2}, \dots, C_{5b}$ and there are at the most $2b$ consecutive remaining components and will be a linear combination of atmost two distinct sets of b or less consecutive columns. Therefore no more columns can be added to H .

Thus the $(5b, b)$ binary code which is the null space of the matrix H as constructed above will correct all 2-repeated bursts of length b or less.

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